

Alexander Woo

awoo@uidaho.edu

(208)-885-6741 (office)

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The Halting Problem

Thm: There cannot a computer program that takes as input a computer program (as a long string in C) and outputs whether this program will halt or not given that input.

and input to that program

Pf: Suppose for contradiction there was such a program, packaged as a C function

```
boolean haltp(char* program,  
              char* input)
```

...

end

Write a new program:

```
void breaker(char x)
  if haltp(x, x)
    go into infinite loop
  else
    terminate
```

Call
↓
breaker (this string)

This is a contradiction, since breaker breaks itself. if it terminates, ~~it won't~~ haltp returns true and it loops; if it loops, haltp returns false and it terminates.

~~There~~ Conclusion: there is no such program.

A false theorem:

"Theorem": Every positive integer can be described in less than 20 words.

"Pf": Suppose not - then there exists a smallest pos. int. not describable in ≤ 20 words. We can describe this integer as "The smallest pos. int. not describable in less than 20 words" \square

What's wrong is that we haven't been precise about what "describe" means.

This throws into question the previous argument

Goal for semester: Give a mathematical formalization of "computer program," "input," "output," "feeding a program as input to another program" - so that we can be sure this argument works.

Outline for semester:

singular: automaton
plural: automata

- 1 We make mathematical def'n's of several kinds of automata - formal mathematical models of computers w/varying levels of complexity and capabilities.
- 2 Automata will take a string as input and decide "true" or "false" for the input.
- 3 We can think about the set of strings that an automaton says "true" to. This set is a language. We will study ~~how~~ what properties a language must have to be accepted by a particular type of automaton.

E.g.: { strings of a's and b's where the number of a's is divisible by the number of b's }

abaa is in this language (3 is div by 1)

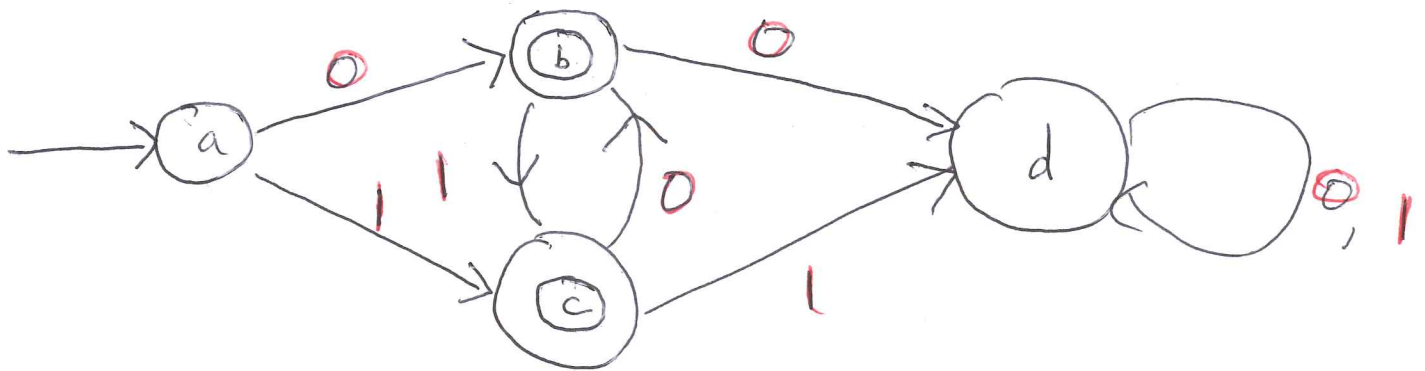
abaab is not (3 is not div by 2)

E.g. { strings of ('s and) 's
~~where~~ which form proper strings of
parens. - i.e. $\#('s = \#) 's$, and at
any point, we have seen at least as many
('s as) 's. }

4 Think about a language in terms of rules
for generating all possible strings in the
language. This collection of rules is a
grammar.

Study how type of automaton
↑
restrictions on language
↓
type of grammar

Deterministic Finite Automata (DFA)



0101110
a b c b c d d d → not an accept state, so false

A DFA has a finite number of states

One is the initial state (a)

When fed a letter, it ~~changes state~~ goes to some state depending on the current state and the letter it is fed.

Feeding a string just means feeding letters one at a time. If, at the end of the string, we are in a final state (or accept state), it says "true," and otherwise "false"

We just worked out - this DFA accepts exactly the strings which strictly alternate 0's and 1's.