

Getting a regular expression out of an NFA/DFA

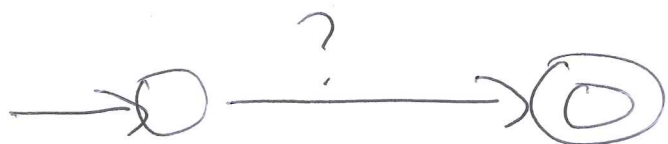
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A super-NFA is an NFA where we allow arbitrary regular expressions to label the transitions.

One way of thinking of an NFA - for a transition, we have ~~to~~ to pay a toll out ^{given by the label} of the beginning of the string to use it - and a string is accepted if we can use up our string and get to a final state.

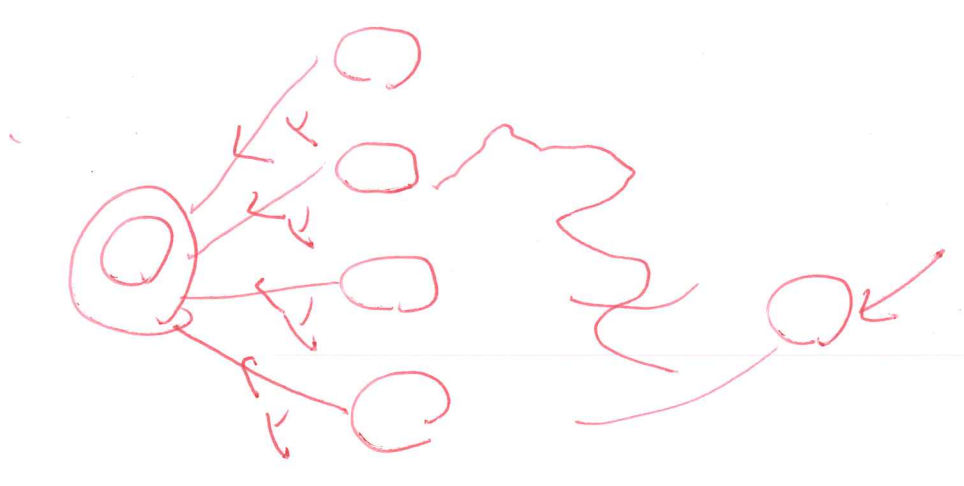
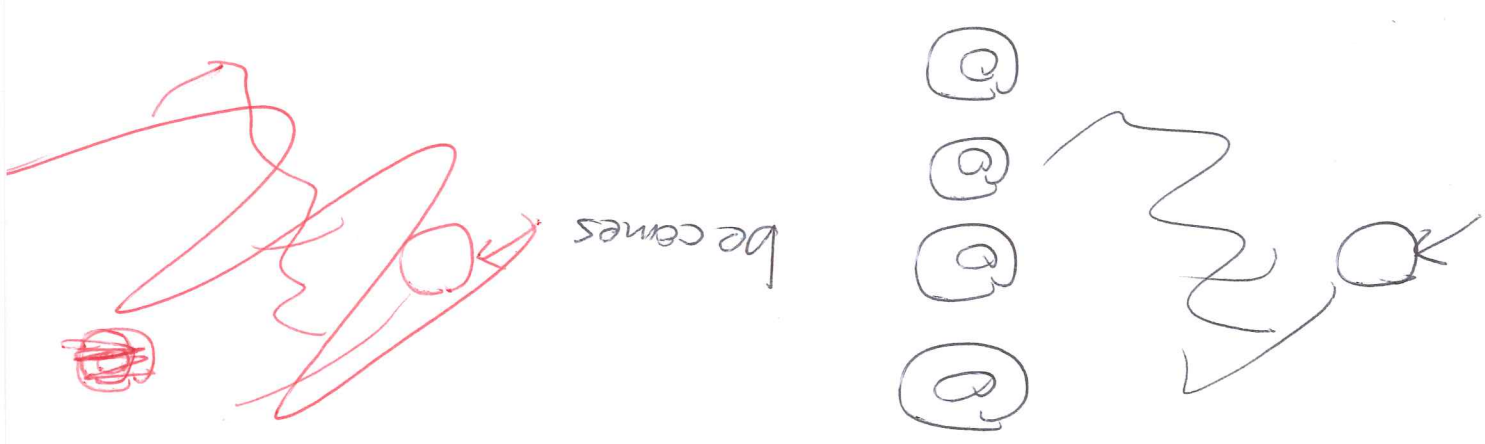
For a super-NFA, I can pay any toll that matches the r.e. on a transition.

Getting a r.e. out of an NFA is the same as converting an NFA to an equivalent super NFA of the form

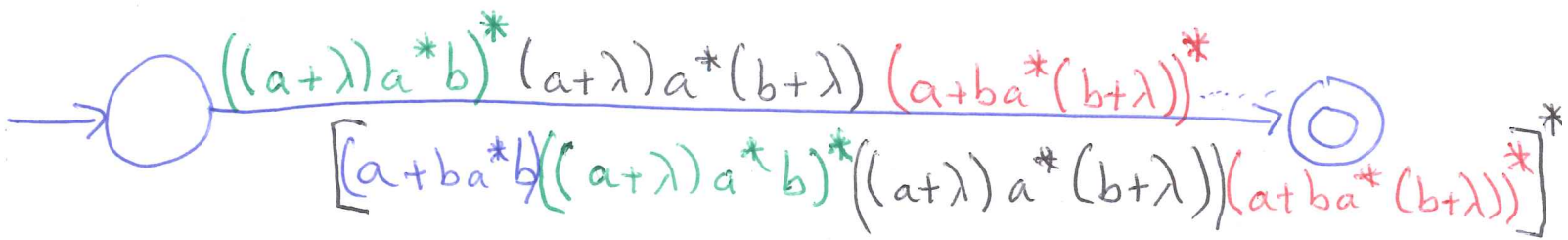


So we need to figure out how to get rid of states (in the process making transition labels more complicated) while keeping the accepted language the same.

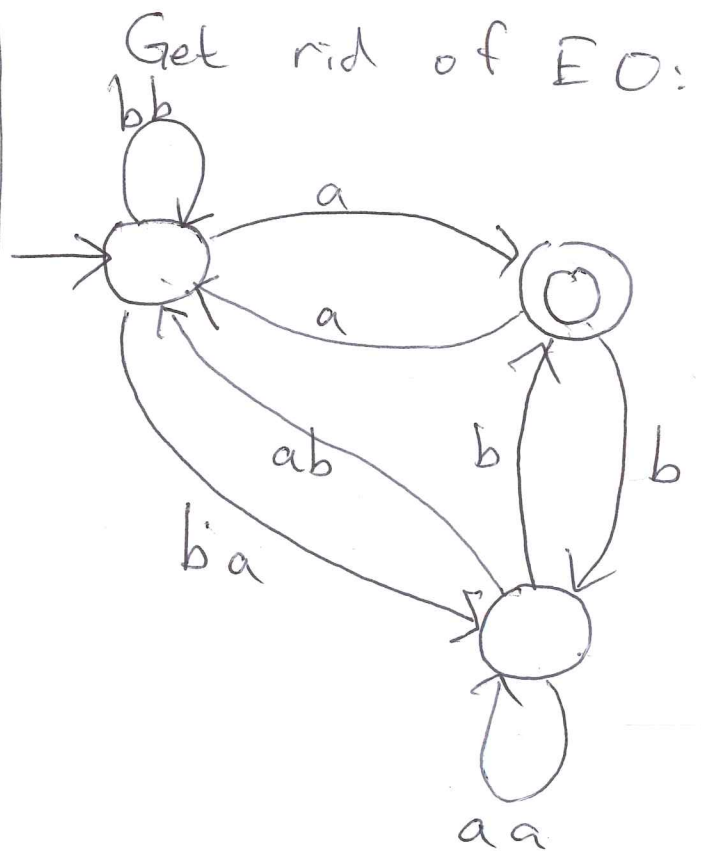
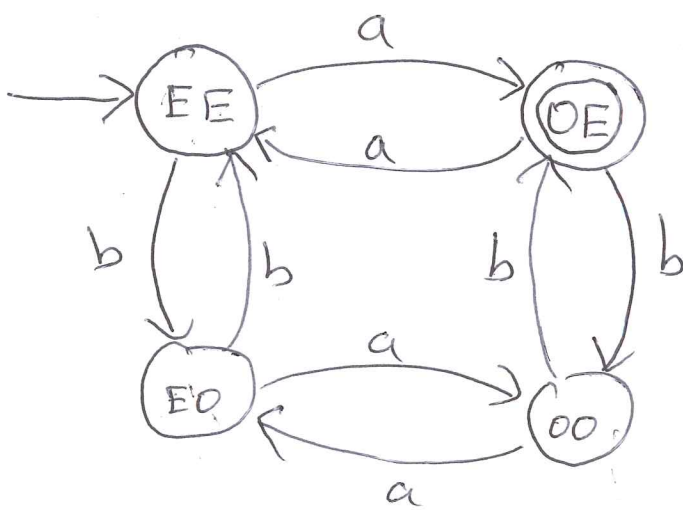
0) Change my NFA to have only one final state.



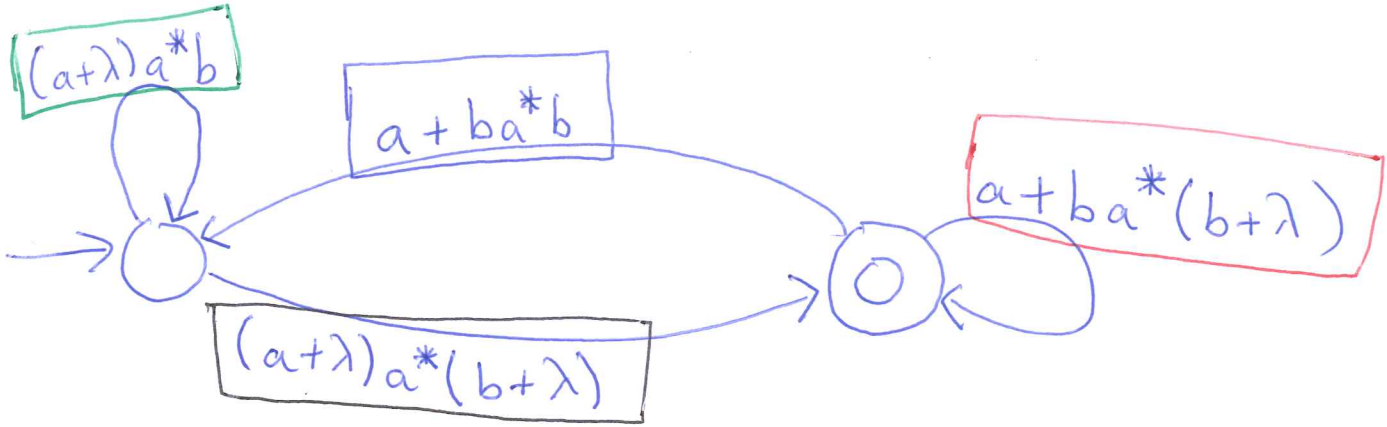
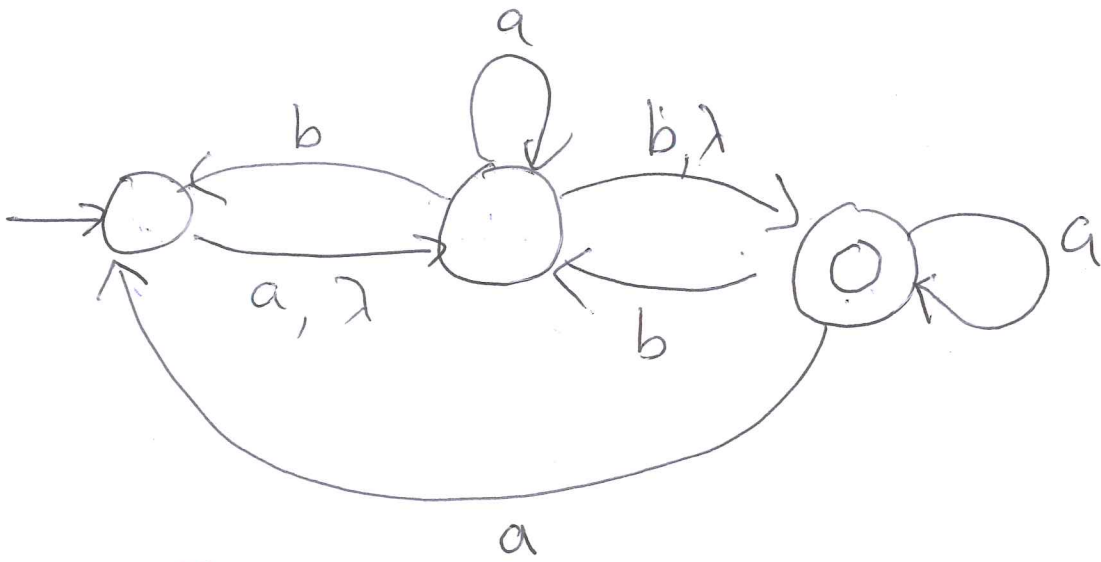
When we are down to one final and one initial state:



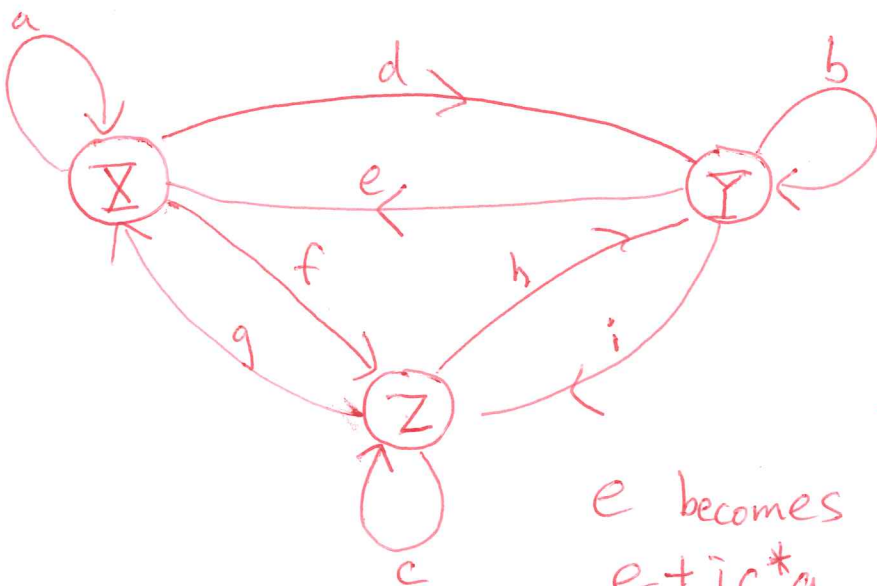
Find a regular expression for the language of strings with an odd # of a's and an even # of b's:



1) Get rid of non-final states as follows:



Formula

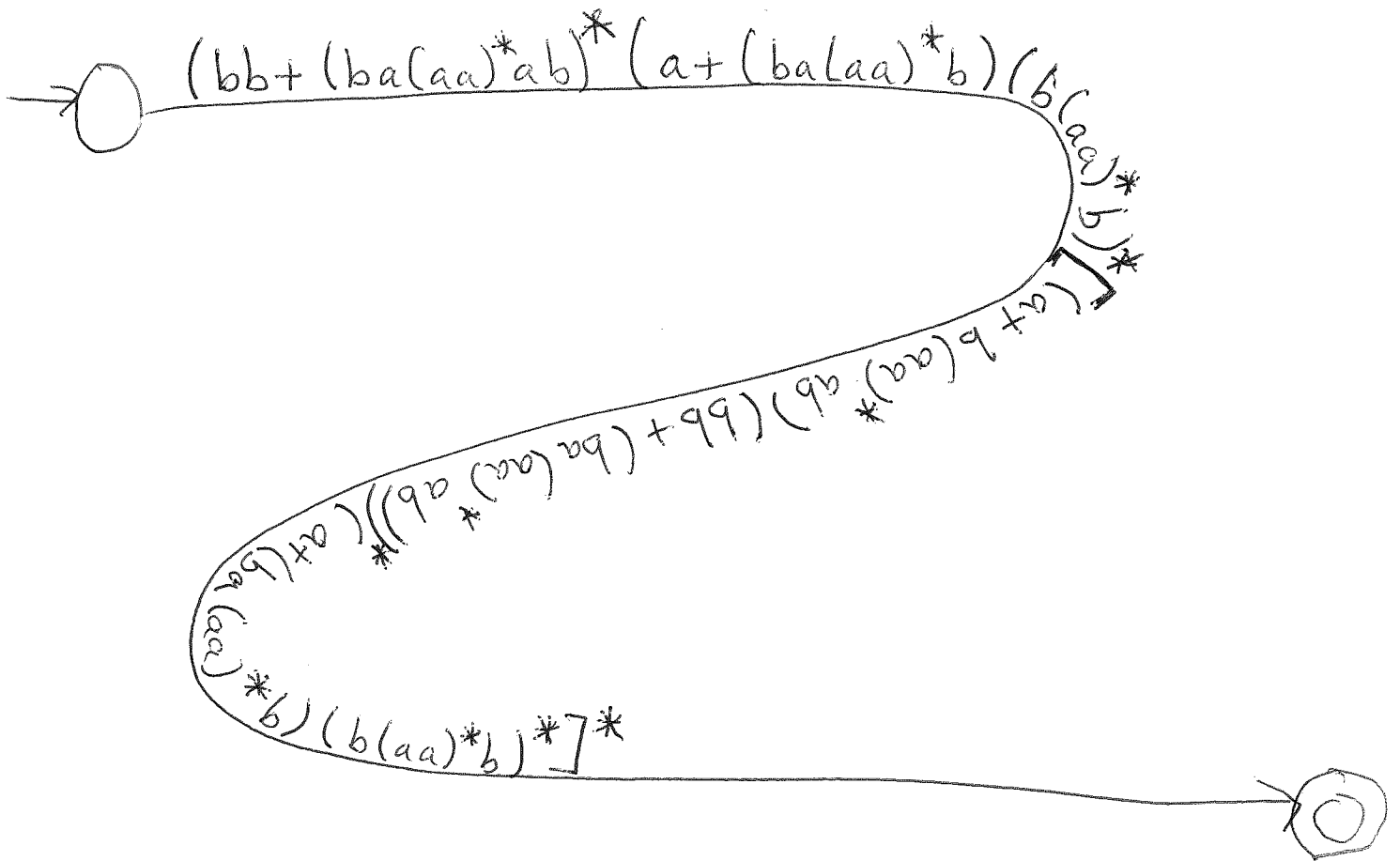
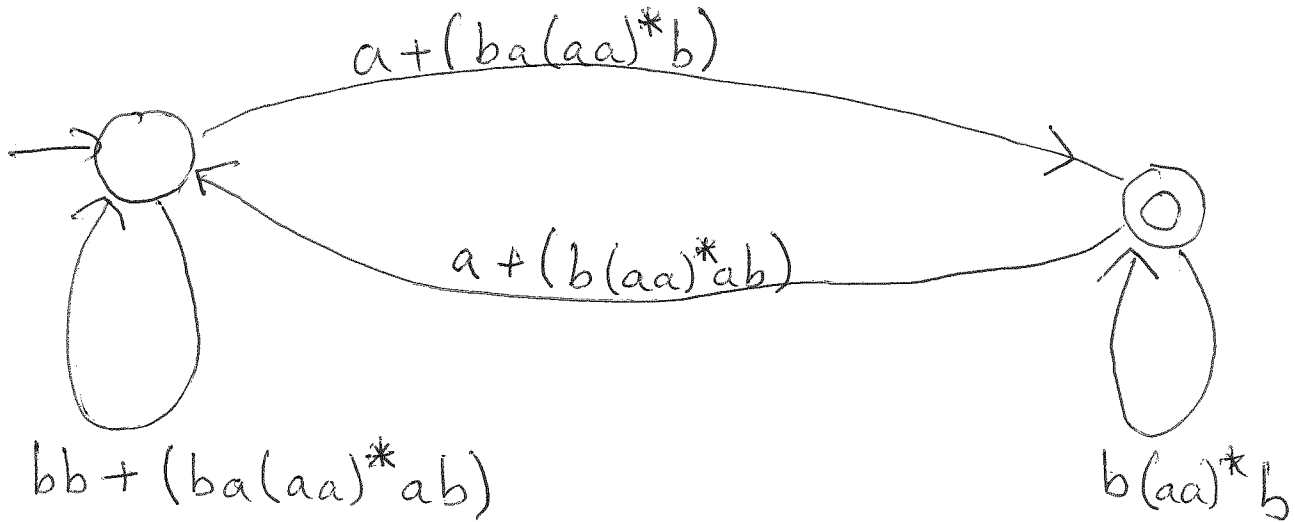


e becomes $e+ic^*g$
 b becomes $b+ic^*h$

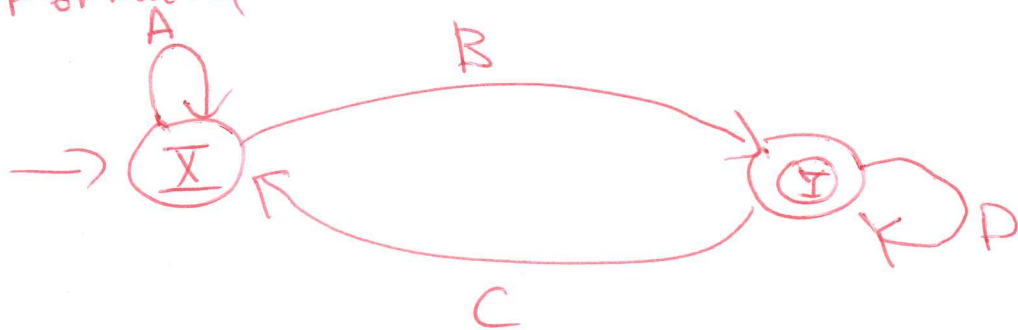
To get rid of Z,
 between any
 states X and Y

~~a~~ a becomes $a+fc^*g$
 d becomes $d+fc^*h$

Get rid the non-init, non-final state.



Formula



is equiv to:

$$A^*BD^*(CA^*BD^*)^*$$

