

Useful ways to turn CFGs into equivalent CFGs

10/5/18

Goal: Monday we'll talk about "normal forms" certain restrictions on what our productions look like - we want to convert every CFG into a CFG in normal form.

To do this, we need to know ways to convert CFGs into equivalent CFGs.

Useless variables:

E.g. $S \rightarrow \cancel{xAy} \mid yBy$

$$A \rightarrow xSx \mid Bx$$

$$B \rightarrow ySx \mid Ay \mid \lambda$$

$$C \rightarrow xBy \quad \text{useless}$$

How do we find these useless variables?

~~start at S~~

$$\text{useful} = (S)$$

take a useful variable,
add every var it produces
into the list of useful var.

2 lists of variables:
analyzed
useful

Put the just analyzed variable
~~it~~ into the analyzed list

repeat until every useful variable
is analyzed.

E.g. $S \rightarrow xAy \mid yBy$

$$A \rightarrow xSx \mid Bx \mid \lambda$$

$$B \rightarrow \del{Sx} \del{By} Cy \text{] both useless}$$

$$C \rightarrow Bx$$

Here, once you get a B, you can
never get rid of it (more generally,
you cannot get rid of all vars)

So B is useless - you can't produce
an actual string of letters from it.

How to check?

Every production ending in only letters
makes the var good.

Every prod. ending in only letters and
good vars makes the LHS var good.

If going through the list of prods
doesn't expand list of good ~~letter~~ vars,
you are done.

Step 0 for simplifying grammars:

Get rid of all useless variables.

Getting rid of λ -productions:

E.g. $S \rightarrow xAy \mid yBy \mid yy$

$$A \rightarrow xSx \mid Bx \mid x$$

$$B \rightarrow ySx \mid Ay \mid \lambda$$

Add ~~pro~~ all the productions that do

$z \rightarrow zBz \rightarrow \dots$
as one step productions

~~Since~~ If I have a derivation

$$S \rightarrow xAy \rightarrow xBxy \rightarrow xxy$$

Change the grammar so that I have

$$S \rightarrow xAy \rightarrow xxy$$

General principle:

If I have a production

$$A \rightarrow w \quad (A \text{ a var, } w \text{ a string of vars \& letters})$$

I can get rid of the production at the price of adding an extra production to every production that produces A.

Careful

$A \rightarrow x B_x B_y \mid x x B_y \mid x B_x y \mid x x y$
 $B \rightarrow \lambda \mid \dots$

all the possibilities for subbing λ for B

Getting rid of $A \rightarrow B$ productions
(unit productions)

E.g. $S \rightarrow x A_y \mid B_x$

to get rid of this, add

$A \rightarrow \lambda \mid x S \mid x A_x \mid B_y \mid S_y \mid \lambda$

$B \rightarrow x A_x \mid B_y \mid S_y \mid \lambda$

Careful - we can have a chain of unit productions

$S \rightarrow x A_y \mid B_x$
 $A \rightarrow B \mid x S$
 $B \rightarrow x A_x \mid C \mid \lambda$
 $C \rightarrow B_y \mid A$

$A, B, \text{ \& } C$
 can all go
 to each
 other,
 so all prods
 of any of
 them are prods
 of all.

$$S \rightarrow xAy \mid \bar{x}$$

$$\bar{x} \rightarrow xS \mid xAx \mid \lambda \mid \bar{y}$$

~~S~~ $\bar{x} = \bar{B}$ ← A on top of B on top of C.

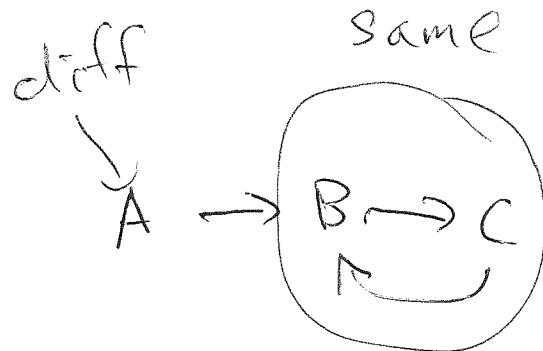
or a non loop:

$$S \rightarrow xAy \mid Bx$$

$$A \rightarrow B \mid xS$$

$$B \rightarrow xAx \mid C \mid \lambda$$

$$C \rightarrow By \mid B$$



$$S \rightarrow xAy \mid \bar{x}$$

$$A \rightarrow xS \mid xAx \mid \lambda \mid \bar{y}$$

$$\bar{x} \rightarrow xAx \mid \lambda \mid \bar{y}$$

merge all
prods of
B + C,
add things
B/C produce
as productions
of A.